

Numerical Optimization

Lecture Notes #14

Practical Newton Methods — Hessian Modifications

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Outline

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 - Robust Inexact Newton Methods
- 2 Hessian Modifications
 - Eigenvalue Modification
 - $\mathbf{B} = \mathbf{A} + \tau \mathbf{I}$
 - Gershgorin Modification



Quick Recap: Building Robust Inexact Newton Methods

We looked at combining a modified version of the linear CG-solver (or preferably a PCG(M)-solver) with a line-search algorithm to produce an almost “unbreakable” approximate Newton method.

The modification to the CG-solver comprise of an additional termination criterion for the case where the local Hessian ($\nabla^2 f(\bar{\mathbf{x}}_k)$) is not positive definite, and we get a CG-internal search direction for which $\bar{\mathbf{p}}^T \nabla^2 f(\bar{\mathbf{x}}_k) \bar{\mathbf{p}} \leq 0$, i.e the search takes into a part of space with negative curvature.

The worst we do (in a particular iteration) is to take a steepest descent step.

Potential Outstanding Problem: $\bar{\mathbf{p}}^T \nabla^2 f(\bar{\mathbf{x}}_k) \bar{\mathbf{p}}$ small and positive \rightsquigarrow long step.



Quick Recap: Building Robust Inexact Newton Methods

We also discussed how to specify the **forcing sequence** $\{\eta^{(k)}\}$ for the tolerance termination criterion ($\|\bar{\mathbf{r}}_k\| \leq \eta^{(k)} \|\nabla f(\bar{\mathbf{x}}_k)\|$) so that the overall convergence rate of the resulting algorithm is quadratic (when $B_k = \nabla^2 f(x_k)$) or super-linear (when $B_k \approx \nabla^2 f(x_k)$).

We also hinted at a different approach to dealing with non-positive definite Hessians in the direct-linear-solver-framework — a modification of the Hessian ($\nabla^2 f(\bar{\mathbf{x}}_k) + E_k$) so that the resulting matrix is sufficiently positive definite; today we take a closer look at this approach.



We look at modifying the Hessian matrix $\nabla^2 f(\bar{\mathbf{x}}_k)$ by either explicitly or implicitly adding a matrix E_k (usually a multiple of the identity matrix) so that the resulting matrix

$$B_k = \nabla^2 f(\bar{\mathbf{x}}_k) + E_k$$

is **sufficiently positive definite** (all the eigenvalues of B_k are bounded away from zero.)

There are a number of different approaches, we look at a few...

- Eigenvalue Modification
- Direct and Indirect modification of the Hessian



Since $\nabla^2 f(\bar{\mathbf{x}}_k)$ is symmetric we can always find an orthonormal matrix Q_k and a diagonal matrix $\Lambda_k = \text{diag}(\lambda_1, \lambda_2, \dots, \lambda_n)$ so that (dropping the subscripts k)

$$\nabla^2 f(\bar{\mathbf{x}}) = Q\Lambda Q^T = \sum_{i=1}^n \lambda_i \bar{\mathbf{q}}_i \bar{\mathbf{q}}_i^T.$$

For simplicity of argument, let us assume $Q = I$ (we can get to this scenario by an appropriate change of variables.)

Example:

$$\nabla f(\bar{\mathbf{x}}) = \begin{bmatrix} 1 \\ -3 \\ 2 \end{bmatrix}, \quad \nabla^2 f(\bar{\mathbf{x}}) = \text{diag}(10, 3, -1) \Rightarrow \bar{\mathbf{p}}^N = \begin{bmatrix} -0.1 \\ 1 \\ 2 \end{bmatrix}$$

and $\nabla f(\bar{\mathbf{x}})^T \bar{\mathbf{p}}^N = 0.90$, hence $\bar{\mathbf{p}}^N$ is **not a descent direction**.
(continued...)



Idea#1: Replace negative eigenvalues by some positive number δ , e.g.
 $\delta = \sqrt{\epsilon^{\text{mach}}}$

In 32-bit double precision (and Matlab) $\epsilon^{\text{mach}} \approx 10^{-16}$, so $\delta = 10^{-8}$ seems like a reasonable choice(?) We can express the Hessian modification as

$$B_k = \sum_{i=1}^2 \lambda_i \bar{\mathbf{q}}_i \bar{\mathbf{q}}_i^T + \delta \bar{\mathbf{q}}_3 \bar{\mathbf{q}}_3^T \quad \left[= \sum_{i=1}^n \max(\lambda_i, \delta) \bar{\mathbf{q}}_i \bar{\mathbf{q}}_i^T \right]$$

We now have

$$B_k = \text{diag}(10, 3, 10^{-8}) \Rightarrow \bar{\mathbf{p}} \approx \begin{bmatrix} -0.1 \\ 1 \\ -200,000,000 \end{bmatrix}$$

We notice that $\bar{\mathbf{p}}$ is approximately parallel to $\bar{\mathbf{q}}_3$, and **huge**...



The long step length violates the spirit of Newton's method — recall that the quadratic convergence properties come from a **local** argument with the Taylor expansion.

Idea#2: Replace negative eigenvalues by $-\lambda_i$

Now $B_k = \text{diag}(|\lambda_1|, |\lambda_2|, \dots, |\lambda_n|)$, and in our example we get

$$\bar{\mathbf{p}} = \begin{bmatrix} -0.1 \\ 1 \\ -2 \end{bmatrix}, \quad \nabla f(\bar{\mathbf{x}})^T \bar{\mathbf{p}} = -7.1, \quad \text{descent direction!}$$

This seems to work?!?

It may reorder the eigenvalues (and thus the "importance" / ordering of subspaces), i.e.

$$\lambda_1 < \lambda_2 < \lambda_3, \quad \text{but} \quad |\lambda_2| < |\lambda_1| < |\lambda_3|.$$



Eigenvalue Modification

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Let's reconsider Idea#1, what went wrong? When we solve $B\bar{\mathbf{p}} = -\nabla f(\bar{\mathbf{x}})$ we get

$$\bar{\mathbf{p}} = -B^{-1}\nabla f(\bar{\mathbf{x}}) = -\sum_{i=1}^2 \frac{1}{\lambda_i} \bar{\mathbf{q}}_i (\bar{\mathbf{q}}_i^T \nabla f(\bar{\mathbf{x}})) - \frac{1}{\delta} \bar{\mathbf{q}}_3 (\bar{\mathbf{q}}_3^T \nabla f(\bar{\mathbf{x}})),$$

it's clearly the right-most term that makes us violate the spirit of Newton's method.

We could simply just drop this term (*i.e.* ignore the subspace corresponding to negative eigenvalues), or

Select δ so that we ensure that the step length is not excessive (**trust-region** flavor!).

Bad news: There is no accepted "best" way of modifying the Hessian in this manner.



Eigenvalue Modification

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If we for a moment "forget" about the issue of selecting δ so that the step length is reasonable, we can ask the question **"what is the smallest change to A , which gives us an positive definite matrix B ?"**

The answer depends on how we measure... Two standard measures are the **Frobenius norm** $\|A\|_F$, and the **Euclidean norm** $\|A\|$

$$\|A\|_F^2 = \sum_{i,j} a_{ij}^2, \quad \|A\| = \max_{\|\bar{\mathbf{x}}\|=1} \bar{\mathbf{x}}^T A \bar{\mathbf{x}} = \max |\text{eig}(A)|.$$

If we use the Frobenius norm, the smallest change is of the type **"change negative eigenvalues to small positive ones:"**

$$B = A + \Delta A, \quad \text{where } \Delta A = Q \text{diag}(\tau_i) Q^T, \quad \tau_i = \begin{cases} 0 & \lambda_i \geq \delta \\ \delta - \lambda_i & \lambda_i < \delta \end{cases}$$



Eigenvalue Modification

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If, on the other hand, we use the Euclidean norm the smallest change includes a multiple of the identity matrix, *i.e.* **"shift the eigenvalue spectrum, so all eigenvalues are positive:"**

$$B = A + \Delta A, \quad \text{where } \Delta A = \tau I, \quad \tau = \max(0, \delta - \lambda_{\min}(A))$$

We recognize this type of modification to A from our discussion on **"Nearly exact solutions to the subproblem"** for trust-region methods (Lecture #9)...

Both constant-diagonal — τI — and "Frobenius-style" — $Q \text{diag}(\tau_i) Q^T$ — modifications are used in production software. Generally they do not rely on an exact spectral decomposition (full computation of the eigenvalues) of the Hessian, but use a cousin of Gaussian Elimination (usually the Cholesky factorization) which allows introduction of modifications indirectly.

 $B = A + \tau I$

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In adding a multiple of the identity matrix, we would like to identify a scalar τ so that

$$\tau = \max\left(0, \delta - \lambda_{\min}(A)\right).$$

Usually we do not have access to $\lambda_{\min}(A)$, so we have to use some clever heuristic to get an estimate and generate

$$\begin{cases} \tau = 0 & \text{if } \lambda_{\min}(A) \geq \delta \\ \tau \geq \delta - \lambda_{\min}(A) & \text{if } \lambda_{\min}(A) < \delta \end{cases}$$

It is important not to select a value of τ that is unnecessarily large, since this biases the direction toward the steepest descent direction.



$B = A + \tau I$

The following algorithm uses the fact that

$$|\lambda_i| \leq \|A\|_F, \quad \forall i = 1, 2, \dots, n$$

it is quite expensive since a new factorization is attempted in each loop, further the generated τ may be unnecessarily large.

Algorithm

```

 $\beta = \|A\|_F$ ,  $k=0$ 
if(  $\min(a_{ii}) > 0$  ) {  $\tau_0 = 0$  } else {  $\tau_0 = \beta/2$  } endif
while(  $k < \text{maxiter}$  )
  ATTEMPT (Incomplete) Cholesky Factorization
   $LL^T = A + \tau_k I$ 
  if(  $\text{successful\_factorization}$  ), return(L)
  else,  $\tau_{k+1} = \max(2\tau_k, \beta/2)$ 
  endif
end(while)
    
```



$B = A + \text{diag}(\bar{d}^{\text{add}})$ — Breaking Cholesky

It is more efficient to let the Cholesky factorization routine directly modify the matrix A so that the factorization succeeds.

What can go wrong in Cholesky factorization?

We look at the Cholesky factorization in LDL^T -form — set $M = LD^{1/2}$ to get to MM^T form.

Algorithm: Cholesky Factorization, LDL^T -form

```

for j = 1:n
   $c_{jj} = a_{jj} - \sum_{s=1}^{j-1} d_s l_{js}^2$ 
   $d_j = c_{jj}$  --- The diagonal entries in  $D$  (must be  $\geq \delta$ )
  for i = (j+1):n
     $c_{ij} = a_{ij} - \sum_{s=1}^{j-1} d_s l_{is} l_{js}$ 
     $l_{ij} = c_{ij}/d_j$  --- We don't want  $l_{ij}$  to be too large
  end
end
    
```



$B = A + \text{diag}(\bar{d}^{\text{add}})$ — Modifying Cholesky

If we want to require that the matrix LDL^T is sufficiently positive definite, we simply modify the elements d_j :

$$d_j = c_{jj} \rightarrow d_j = \max(c_{jj}, \delta)$$

Usually, we also want to have a bound on the size of the off-diagonal entries of $M = LD^{1/2}$, i.e. $|m_{ij}| \leq \beta$ ($i > j$), we set

$$\theta_j = \max_{j < i \leq n} |c_{ij}|$$

and let

$$d_j = c_{jj} \rightarrow d_j = \max\left(c_{jj}, \delta, \left[\frac{\theta_j}{\beta}\right]^2\right)$$

we have

$$|m_{ij}| = |l_{ij} \sqrt{d_j}| = \frac{|c_{ij}|}{\sqrt{d_j}} \leq \frac{|c_{ij}| \beta}{\theta_j} \leq \beta.$$



$B = A + \text{diag}(\bar{d}^{\text{add}})$ — Modifying Cholesky

Finally, we throw in an absolute value on the c_{jj} term for good measure, and come up with

$$d_j = \max\left(|c_{jj}|, \delta, \left[\frac{\theta_j}{\beta}\right]^2\right), \quad d_j^{\text{add}} = d_j - c_{jj}$$

This exactly what the module `choldecomp()` in the old default project does! (With some modifications for computational efficiency — the algorithm generates the factorization directly in LL^T -form)

Old Default Project

choldecomp()	
Implementation	Theory (here)
maxoff1	β
min1	$\sqrt{\delta}$
maxadd	$\max(\text{diag}(\bar{d}^{\text{add}}))$



Gershgorin Modification

choldecomp() and modelhess()

Theorem (Gershgorin's circle theorem)

tells us where the eigenvalues of a matrix are located:

$$|\lambda_i - a_{ii}| \leq \sum_{j \neq i} |a_{ij}|, \quad i = 1, \dots, n.$$

Now given a matrix A , let \mathbf{b}_1 be the smallest value which makes $A + \mathbf{b}_1 I$ positive definite from the Gershgorin circle theorem.

Let $\mathbf{b}_2 = \text{maxadd}$ from choldecomp(), and let $\mu = \min(\mathbf{b}_1, \mathbf{b}_2)$. Now, $A + \mu I$ is guaranteed to be positive definite.

Old Default Project

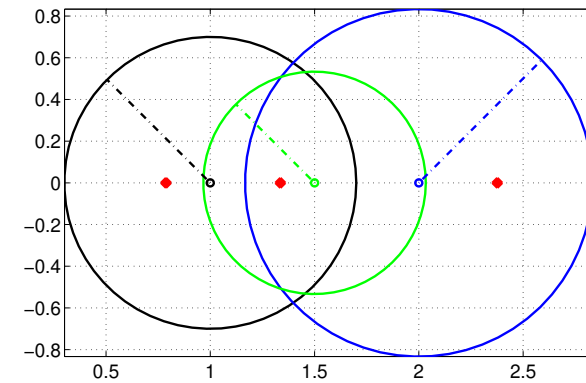
This is essentially modelhess(). In addition modelhess() returns the LL^T -decomposition of $A + \mu I$, and there are tests prior to the first call to choldecomp() which takes care of negative diagonal elements of A and large off-diagonal elements of A .

Note that modelhess() is similar to the algorithm on slide #13, but requires at most two calls to a Cholesky factorization algorithm.



Gershgorin's Circle Theorem: Illustration

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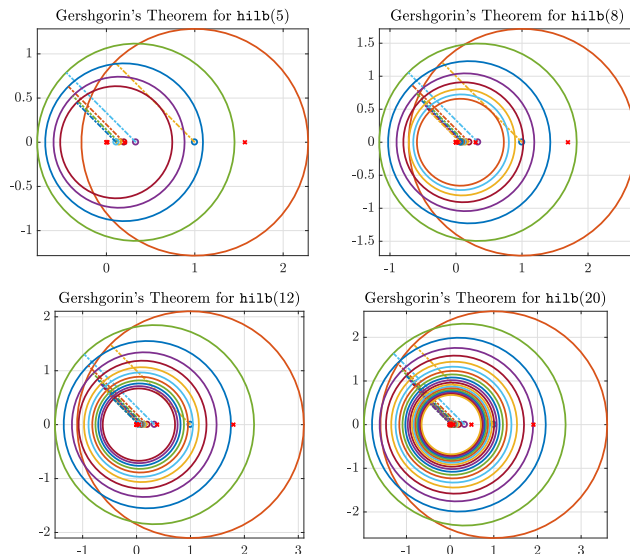


$$A = \begin{bmatrix} 1 & 1/2 & 1/5 \\ 1/2 & 2 & 1/3 \\ 1/5 & 1/3 & 3/2 \end{bmatrix}, \quad \lambda(A) = \{0.7875, 1.3363, 2.3762\}$$



Gershgorin's Circle Theorem: Illustration

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Project Expectation and Deliverables

Clarified

- Solve a larger optimization problem (see e.g. the “examples of past projects” handout from last time).
- You can look at different types of methods; performance for different test functions, etc... **BEST**: something relevant to your thesis project.
- **Deliverables**:
 - *Project Proposal* — 1 page, Due 11/16/2018
 - *Presentation* — 12–15 minutes, in-class (starting 12/10/2018)
 - *email* — presentation + code(s). (after presentation)



Next...

- Practical Newton Methods: Trust-Region Newton Methods
- Calculating Derivatives: Finite Differencing & Automatic Differentiation
- Quasi-Newton Methods...



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