

— (3/32)

We looked at stability for multistep schemes. — First, we did a complete analysis of the stability picture for the leapfrog scheme,

$$\frac{v_m^{n+1} - v_m^{n-1}}{2k} + a \frac{v_{m+1}^n - v_{m-1}^n}{2h} = 0$$

in which we found bounds for the roots of

$$g(\theta)^2 + \left[2ia\lambda\sin(\theta)\right]g(\theta) - 1 = 0$$

so that  $|g_{\pm}( heta)| \leq 1$  for simple roots and  $|g_{\pm}( heta)| < 1$  for multiple roots.

The analysis for general multi-step scheme has the same "flavor," but we postponed the development of a unified framework for that analysis until today. Last time, we boldly stated that the scheme

$$\frac{3v_m^{n+1} - 4v_m^n + v_m^{n-1}}{2k} + a\frac{v_{m+1}^{n+1} - v_{m-1}^{n+1}}{2h} = f_m^{n+1}$$

with amplification polynomial

$$\Phi(g,\theta) = \left[rac{3+2ia\lambda\sin(\theta)}{2}
ight]g^2 - 2g + rac{1}{2}$$

is unconditionally stable, and order-(2,2) accurate.

Whereas pictures are not proof, the plots of the roots for various values of  $a\lambda$  and  $\theta \in [-\pi, \pi]$  shown on slide 7 seem to indicate that the stability is indeed unconditional.





(12/32)



Definitions and Theorems Examples: Revisited with Theoretical Toolbox in Hand... Algorithm for von Neumann / Schur Polynomials

The Polynomial  $\varphi^*(z)$ 

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(19/32)

For a polynomial of exact degree d, we define the polynomial

$$\varphi^*(z) = \sum_{\ell=0}^d \overline{a}_{d-\ell} z^\ell \equiv \overline{\varphi(1/\overline{z})} z^d,$$

where  $\overline{z}$  is the complex conjugate of z.

We recursively define the polynomial  $\varphi_{d-1}$  of exact degree d-1 by

$$\varphi_{d-1}(z) = \frac{\varphi_d^*(0)\varphi_d(z) - \varphi_d(0)\varphi_d^*(z)}{z} \equiv \frac{\overline{a}_d\varphi_d(z) - a_0\varphi_d^*(z)}{z}$$

We are now ready to state theorems which provide tests for Schur and simple von Neumann polynomials.



# Theorem (von Neumann Polynomial Test)

 $\varphi_d$  is a von Neumann polynomial of degree d, if and only if either

- (a)  $|\varphi_d(0)| < |\varphi_d^*(0)|$  and  $\varphi_{d-1}$  is a von Neumann polynomial of degree d-1, or
- **(b)**  $\varphi_{d-1}$  is identically zero and  $\varphi'_d$  is a von Neumann polynomial.

### Theorem (Conservative Polynomial Test)

 $\varphi_d$  is a conservative polynomial if and only if  $\varphi_{d-1}$  is identically zero and  $\varphi'_d$  is a von Neumann polynomial.

# Theorem (Simple Conservative Polynomial Test)

 $\varphi_d$  is a simple conservative polynomial if and only if  $\varphi_{d-1}$  is identically zero and  $\varphi'_d$  is a Schur polynomial.

# Schur and von Neumann Polynomials

Definitions and Theorems Examples: Revisited with Theoretical Toolbox in Hand... Algorithm for von Neumann / Schur Polynomials

## Building the Toolbox

Polynomial Tests

# Theorem (Schur Polynomial Test)

 $\varphi_d$  is a Schur polynomial of exact degree d if and only if  $\varphi_{d-1}$  is a Schur polynomial of exact degree d-1 and  $|\varphi_d(0)| < |\varphi_d^*(0)|$ .

# Theorem (Simple von Neumann Polynomial Test)

 $\varphi_d$  is a simple von Neumann polynomial if and only if either

- (a)  $|\varphi_d(0)| < |\varphi_d^*(0)|$  and  $\varphi_{d-1}$  is a simple von Neumann polynomial, or
- **(b)**  $\varphi_{d-1}$  is identically zero and  $\varphi'_d$  is a Schur polynomial.

The (somewhat lengthy) proofs, which depend on **Rouché's theorem** (complex analysis) are in Strikwerda pp. 110-114.



1 of 2

Schur and von Neumann Polynomials

Examples: Revisited with Theoretical Toolbox in Hand... Algorithm for von Neumann / Schur Polynomials

### Example #1: Revisited

The scheme had the amplification polynomial

$$arphi_2(z) = \left[rac{3+2ia\lambda\sin( heta)}{2}
ight]z^2 - 2z + rac{1}{2}.$$

It is stable exactly when  $\varphi_2(z)$  is a simple von Neumann polynomial. We make repeated use of the simple-von-Neumann-polynomial-test in order to check stability of the scheme.

We first test  $|\varphi_2(\mathbf{0})|^2 = \frac{1}{4} < \frac{1}{4} \left( 3^2 + 4a^2\lambda^2 \sin^2(\theta) \right) = |\varphi_2^*(\mathbf{0})|^2$ , then define, with (c + di) being the coefficient in front of  $z^2$  in  $\varphi_2(z)$ :

$$\varphi_{1}(z) = \frac{1}{z} \left[ (c-di) \left( (c+di)z^{2} - 2z + \frac{1}{2} \right) - \frac{1}{2} \left( (c-di) - 2z + \frac{1}{2}z^{2} \right) \right]$$
$$= \left( d^{2} + c^{2} - \frac{1}{4} \right) z + (1 - 2c + 2id)$$

- (20/32)

2 of 2

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Definitions and Theorems Examples: Revisited with Theoretical Toolbox in Hand... Algorithm for von Neumann / Schur Polynomials

1 of 2

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(24/32)

### Example #2: Revisited

The scheme had the amplification polynomial

$$\varphi_2(z) = \left[7+4i\beta\right]z^2 - \left[8-2i\beta\right]z+1$$

it is stable exactly when  $\varphi_2(z)$  is a simple von Neumann polynomial. We make repeated use of the simple-von-Neumann-polynomial-test in order to check stability of the scheme.

With  $\beta = a\lambda \sin(\theta)$ , we first test  $|\varphi_2^*(\mathbf{0})| = |\mathbf{7} - \mathbf{4i}\beta| > \mathbf{1} = |\varphi_2(\mathbf{0})|$ , then define

$$\begin{split} \varphi_{1}(z) &= \frac{1}{z} \bigg[ (7 - 4i\beta) \left( \big[ 7 + 4i\beta \big] z^{2} - \big[ 8 - 2i\beta \big] z + 1 \right) \\ &- 1 \left( \big[ 7 - 4i\beta \big] - \big[ 8 + 2i\beta \big] z + z^{2} \right) \bigg] \\ &= 4 \left( \left( (12 + 4\beta^{2}) z + ((2\beta^{2} - 12) + 12i\beta) \right). \end{split}$$

 

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 Schur and von Neumann Polynomials
 - (22/32)

 Schur and von Neumann Polynomials
 Definitions and Theorems Examples: Revisited with Theoretical Toolbox in Hand... Algorithm for von Neumann / Schur Polynomials

 Example #3: Revisited
 1 of 4

In this case the amplification polynomial is given by

$$\varphi_3(z) = \left[23 - 12\alpha + 12i\beta\right]z^3 - \left[21 - 12\alpha - 12i\beta\right]z^2 - 3z + 1$$

where

$$\alpha = \frac{a^2 \lambda^2 \sin^2\left(\frac{\theta}{2}\right)}{1 - \frac{2}{3} \sin^2\left(\frac{\theta}{2}\right)} \in [\mathbf{0}, \, \mathbf{3a}^2 \lambda^2], \quad \beta = \frac{a \lambda \sin\left(\theta\right)}{1 - \frac{2}{3} \sin^2\left(\frac{\theta}{2}\right)} \in [-\mathbf{a}\lambda\sqrt{\mathbf{3}}, \, \mathbf{a}\lambda\sqrt{\mathbf{3}}].$$

The first check  $|arphi_3(0)| < |arphi_3^*(0)|$  can be expressed as  $|arphi_3^*(0)|^2 - |arphi_3(z)|^2 > 0$ , and we get

$$|\varphi_3^*(0)|^2 - |\varphi_3(0)|^2 = 24(2-\alpha)(11-6\alpha) + 12^2\beta^2$$

we see that we must require  $0 \le \alpha \le \frac{11}{6}$  for stability.

Now,  $\varphi_1(z)$  is a simple von Neumann polynomial as long as

Stability

Schur and von Neumann Polynomials

$$\left(d^2+c^2-rac{1}{4}
ight)^2 \geq (1-2c)^2+4d^2=1+4c^2+4d^2-4c^2$$

where  $c = \frac{3}{2}$ , and  $d = a\lambda \sin(\theta)$ .

Plugging in we must have

Example #1: Revisited

 $a^{4}\lambda^{4}\sin^{4}\left(\theta\right)+4\,a^{2}\lambda^{2}\sin^{2}\left(\theta\right)+4\geq4\,a^{2}\lambda^{2}\sin^{2}\left(\theta\right)+4$ 

Which holds strictly for  $sin(\theta) \neq 0$ , and with equality when  $sin(\theta) = 0$ .

**Conclusion:** The scheme is unconditionally stable.



$$\varphi_1(z) = 4\left(\left(12 + 4\beta^2\right)z + \left((2\beta^2 - 12) + 12i\beta\right)\right)$$

is a simple von Neumann polynomial if and only if

 $\begin{aligned} |\varphi_1(0)|^2 &= |(2\beta^2 - 12) + 12i\beta|^2 = (12 - 2\beta^2)^2 + 12^2\beta^2 \\ &= \mathbf{144} + \mathbf{96}\beta^2 + \mathbf{4}\beta^4 \le |\varphi_1^*(0)|^2 = (12 + 4\beta^2)^2 = \mathbf{144} + \mathbf{96}\beta^2 + \mathbf{16}\beta^4 \end{aligned}$ 

The inequality holds strictly as long as  $\beta \neq 0$ , in which case we get equality.

Note: Since  $\varphi_1(z)$  only has one root, it is sufficient to bound that root by " $\leq 1$ " in order for  $\varphi_1(z)$  to be a simple von Neumann polynomial.

Conclusion: The scheme is unconditionally stable.

#### Examples: Revisited with Theoretical Toolbox in Hand... Examples: Revisited with Theoretical Toolbox in Hand. Schur and von Neumann Polynomials Schur and von Neumann Polynomials Algorithm for von Neumann / Schur Polynomials Algorithm for von Neumann / Schur Polynomial Example #3: Revisited 2 of 4 Example #3: Revisited 3 of 4 Finally, the polynomial $\varphi_1(z)$ is The polynomial $\varphi_2(z)$ is (after division by the common factor 24) $\varphi_{1}(z) = \left| 120 - 252\alpha + 198\alpha^{2} - 69\alpha^{3} + 9\alpha^{4}(18\alpha^{2} - 69\alpha + 65)\beta^{2} + 9\beta^{4} \right| z$ $\varphi_2(z) = \left| (11 - 6\alpha)(2 - \alpha) + 6\beta^2 \right| z^2$ $+9\beta^{4}+6(5-3\alpha)i\beta^{3}+(3\alpha-5)\beta^{2}-(18\alpha^{3}+102\alpha^{2}+192\alpha-120)i\beta^{3}$ $-2\left|(2-\alpha)(5-3\alpha)-3\beta^2-(11-6\alpha)i\beta\right|z-(2-\alpha-2i\beta),$ $-9\alpha^{4}+69\alpha^{3}-198\alpha^{2}+252\alpha-120$ and The root-condition $|\varphi_1^*(0)|^2 - |\varphi_1(0)|^2 > 0$ translates to $|\varphi_2^*(0)|^2 - |\varphi_2(0)|^2 = 4(5 - 3\alpha) \left| 3(2 - \alpha)^3 + \beta^2(13 - 6\alpha) \right| + 36\beta^4.$ $12\beta^4(5-3\alpha)\left[6\beta^2+(11-6\alpha)(2-\alpha)\right]>0$ This now requires that $0 \le \alpha \le \frac{5}{3} < \frac{11}{6}$ for stability. Ê This holds in the range $0 \le \alpha \le \frac{5}{3}$ ; our strictest bound on $\alpha$ . SAN DIEGO SI UNIVERSIT SAN DIEGO S UNIVERSE Schur and von Neumann Polynomials Peter Blomgren, blomgren.peter@gmail.com Peter Blomgren, blomgren.peter@gmail.com Schur and von Neumann Polynomials - (25/32) · (26/32) Definitions and Theorems Definitions and Theorem Examples: Revisited with Theoretical Toolbox in Hand... Examples: Revisited with Theoretical Toolbox in Hand.. Schur and von Neumann Polynomials Schur and von Neumann Polynomials Algorithm for yon Neumann / Schur Polynomi Algorithm for yon Neumann / Schur Polynon Example #3: Revisited 4 of 4 Example #4: Revisited 1 of 3 We now have that $\varphi_4(z) = z^4 + \frac{4}{3}i\beta\left(2z^3 - z^2 + 2z\right) - 1, \quad \beta = \frac{a\lambda\sin\left(\theta\right)}{1 - \frac{2}{2}\sin^2\left(\frac{\theta}{2}\right)} \in [-a\lambda\sqrt{3}, \ a\lambda\sqrt{3}].$ $\alpha = |a\lambda|^2 \underbrace{\frac{\sin^2\left(\frac{\theta}{2}\right)}{1 - \frac{2}{3}\sin^2\left(\frac{\theta}{2}\right)}}_{3 \le \frac{2}{3} \le \frac{5}{3}$ Here, $|\varphi_4(0)| = |\varphi_4^*(0)| = 1$ . But $\varphi_3(z) \equiv 0$ , hence there is still hope, for $\varphi_4(z)$ being a simple von Neumann polynomial. We must test whether $\psi_3(z) = \frac{3}{4}\varphi'_4(z) = 3z^3 + i\beta(6z^2 - 2z + 2)$ is a **Schur** polynomial. $|\psi_3^*(0)| - |\psi_3(0)| = 3 - |2\beta| > 0$ , as long as $|\beta| < \frac{3}{2}$ . and it follows that the scheme is stable if and only if We form $|\mathbf{a}\lambda| \leq \frac{\sqrt{5}}{2} \approx 0.7454\dots$ $\psi_2(z) = (9 - 4\beta^2)z^2 + (4\beta^2 + 18i\beta)z - 12\beta^2 - 6i\beta$ $|\psi_2^*(0)|^2 - |\psi_2(0)|^2 > 0$ if and only if $(9 - 4\beta^2)^2 > (12\beta^2)^2 + (6\beta)^2$ , which gives $\beta^2 < \frac{9}{64} [\sqrt{41} - 3] < \frac{9}{4}$ . Ê SAN DIEGO S SAN DIEGO S

(27/32)

(28/32)

efinitions and Theoren Examples: Revisited with Theoretical Toolbox in Hand... Algorithm for von Neumann / Schur Polynomials

2 of 3

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(31/32)

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Next, we form

$$\psi_{1}(z) = \left(81 - 108\beta^{2} - 128\beta^{4}\right)z + \left(\left[32\beta^{4} + 144\beta^{2}\right] - i\left[264\beta^{3} - 162\beta\right]\right)z + \left(\left[32\beta^{4} + 16\beta^{2} - 16\beta^{2}\right]z + i\left[32\beta^{4} + 16\beta^{2}\right] - i\left[26\beta^{4} + 16\beta^{2}\right]z + i\left[32\beta^{4} + 16\beta^{2}\right]z$$

The one root is inside the unit circle only if

$$\left(81 - 108\beta^2 - 128\beta^4\right)^2 - \left(\left[32\beta^4 + 144\beta^2\right]^2 + \left[264\beta^3 - 162\beta\right]^2\right) \ge 0.$$

This expression can be factored as

 $3\left(9-4\beta^2\right)\left(3-16\beta^2\right)\left(\underbrace{\beta^2(80\beta^2-72)+81}_{>0}\right)\geq 0.$ 

Hence,  $\psi_1(z)$  is a Schur polynomial for

$$\beta^2 < \frac{3}{16} < \frac{9}{64} \left[ \sqrt{41} - 3 \right].$$

Schur and von Neumann Polynomials

Stability Schur and von Neumann Polynomials

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Definitions and Theorem Examples: Revisited with Theoretical Toolbox in Hand.. Algorithm for von Neumann / Schur Polynomials

# Algorithm for von Neumann / Schur Polynomials

### Algorithm

Start with  $\varphi_d(z)$  of exact degree d, and set NeumannOrder = 0. while (d > 0) do

- 1. Construct  $\varphi_d^*(z)$
- Define  $c_d = |\varphi_d^*(0)|^2 |\varphi_d(0)|^2$ . (\*) 2.
- Construct the polynomial  $\psi(z) = \frac{1}{z} (\varphi_d^*(0) \varphi_d(z) \varphi_d(0) \varphi_d^*(z)).$ 3.
- 4.1. If  $\psi(z) \equiv 0$ , then increase NeumannOrder by 1, and set  $\varphi_{d-1}(z) := \varphi_d'(z).$
- 4.2. Otherwise, if the coefficient of degree d-1 in  $\psi(z)$  is 0, then the polynomial is **not** a von Neumann polynomial of any order, terminate algorithm.
- 4.3. Otherwise, set  $\varphi_{d-1}(z) := \psi(z)$ .

### end-while (decrease d by 1)

(\*) Enforce appropriate conditions on  $c_d$ .

Definitions and Theorem Examples: Revisited with Theoretical Toolbox in Hand... Algorithm for von Neumann / Schur Polynomials

3 of 3

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### Example #4: Revisited

Hence, our final stability condition is

$$|\beta| = rac{|a\lambda\sin(\theta)|}{1-rac{2}{3}\sin^2\left(rac{ heta}{2}
ight)} < rac{\sqrt{3}}{4}.$$

The maximum occurs when  $\cos(\theta) = -1/2$ , and the scheme is stable when  $|\mathbf{a}\lambda| < \frac{1}{4}$ .

Note that even though the scheme is implicit, it is **not** unconditionally stable.



# Comments on the Algorithm

At the end of the algorithm, if the polynomial has not been rejected by 4.2 —

- The polynomial is a von Neumann polynomial of the resulting order (NeumannOrder) provided that all the parameters  $c_d$  satisfy the appropriate inequalities. - These inequalities provide the stability conditions.
- For first-order PDEs, the amplification polynomial must be a von Neumann polynomial of order 1 for the scheme to be stable.
- For second-order PDEs, the amplification polynomial must be a von Neumann polynomial of order 2 for the scheme to be stable.
- Schur polynomials are von Neumann polynomials of order 0.

This analysis can be automated using a symbolic toolbox. — Again, we have reduced something complicated to a deterministic "recipe."

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